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# Multigrid Overview

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Saraniti**

# History of Multi-Grid Method

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- 1964: first paper, Fedorenko, Russia
  - large constants:  $\sim 40,000$  work units, no implementation?
- 1977: Achi Brandt, Israel, made it practical, wrote seminal paper
- late 70's: Nicolaides, Hackbusch, and others proved convergence for certain PDE's; Brandt proved fast convergence
- interest took off around 1981
- but there was (and still is) much skepticism from some because there was little theory
- today used to solve PDE's in many disciplines
- current research: a drive to achieve "textbook efficiency" for general flow simulations (all Mach numbers and Reynolds numbers)
- somewhat superseded by wavelet methods?

## Literature

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Brandt, 1988, The Weizmann Insitute Research in Multilevel Computation: 1988 Report, Proc. Copper Mtn. Conf. on Multigrid Methods, 1989 (53 pp.) *Survey of recent applications. I found this quite thought-provoking.*

Brandt, 1982, Guide to Multigrid Development, in Hackbusch & Trottenberg, eds., Multigrid Methods, pp. 220-312. *Guidelines for multigrid implementers. Long.*

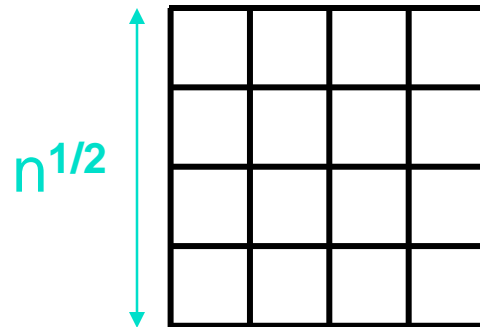
Brandt, 1997, The Gauss Center Research in Multiscale Scientific Computation, Proc. Copper Mtn. Conf. on Multigrid Methods, on web (50 pp.)  
<http://www.wisdom.weizmann.ac.il/research.html> *More esoteric than 1988 report above.*

Brandt, 1980, Multilevel Adaptive Computations in Fluid Dynamics, AIAA J., vol. 18, pp. 1165-1172. *Short, fairly readable.*

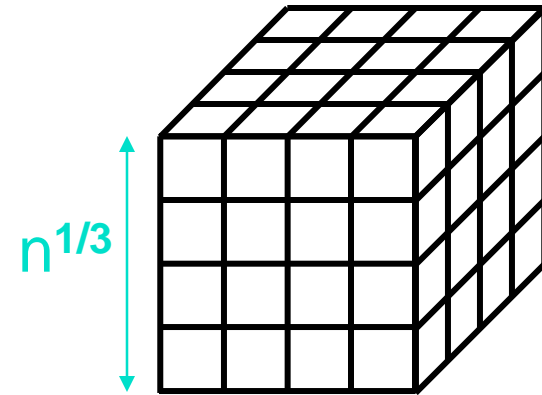
Brandt, 1977, Multi-Level Adaptive Solutions to Boundary-Value Problems, Mathematics of Computation, pp. 333-390. *The seminal paper on multigrid.*

# Complexity of linear solvers

Time to solve  
model problem  
(Poisson's  
equation) on  
regular mesh



2D



3D

<b>Sparse Cholesky:</b>	$O(n^{1.5})$	$O(n^2)$
<b>CG, exact arithmetic:</b>	$O(n^2)$	$O(n^2)$
<b>CG, no precond:</b>	$O(n^{1.5})$	$O(n^{1.33})$
<b>CG, modified IC:</b>	$O(n^{1.25})$	$O(n^{1.17})$
<b>CG, support trees:</b>	$O(n^{1.20}) \rightarrow O(n^{1+})$	$O(n^{1.75}) \rightarrow O(n^{1.31})$
<b>Multigrid:</b>	$O(n)$	$O(n)$

## Outline and Review

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- Review Poisson equation
  - Overview of Methods for Poisson Equation
  - Jacobi's method
  - Red-Black SOR method
  - Conjugate Gradients
  - FFT
- } Reduce to sparse-matrix-vector multiply  
Need them to understand Multigrid
- Multigrid

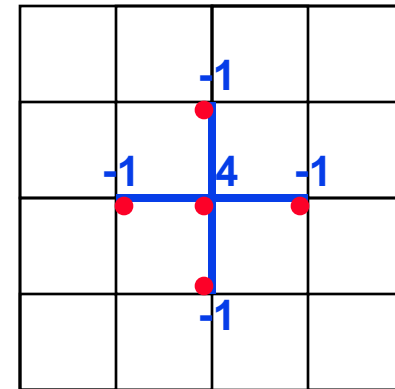
## 2D Poisson's equation

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- ° Similar to the 1D case, but the matrix  $T$  is now

$$T = \begin{pmatrix} 4 & -1 & & -1 & & \\ -1 & 4 & -1 & & -1 & \\ & -1 & 4 & & & -1 \\ -1 & & & 4 & -1 & -1 \\ & -1 & -1 & 4 & -1 & -1 \\ & & -1 & -1 & 4 & -1 \\ & & & -1 & & 4 & -1 \\ & & & & -1 & -1 & 4 & -1 \\ & & & & & -1 & -1 & 4 \end{pmatrix}$$

Graph and “stencil”



- ° 3D is analogous

## Multigrid Motivation

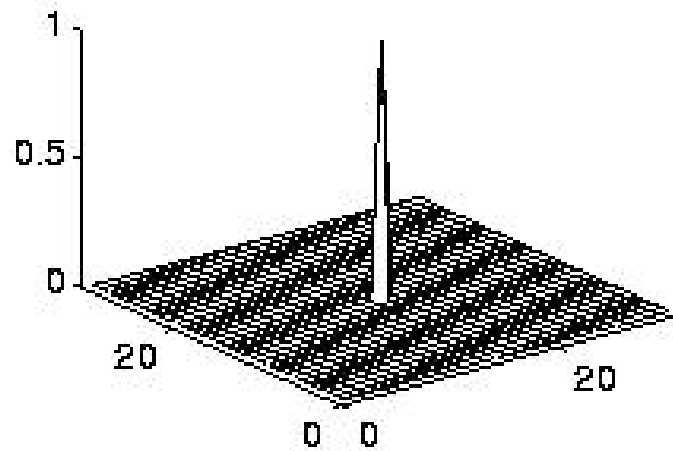
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- **Jacobi, SOR, CG, or any other sparse-matrix-vector-multiply-based algorithm can only move information one grid call at a time for Poisson**
  - New value at each grid point depends only on neighbors
- **Can show that decreasing error by fixed factor  $c < 1$  takes at least  $O(n^{1/2})$  steps**
  - See next slide: true solution for point source like  $\log 1/r$
- **Therefore, converging in  $O(1)$  steps requires moving information across grid faster than just to neighboring grid cell per step**

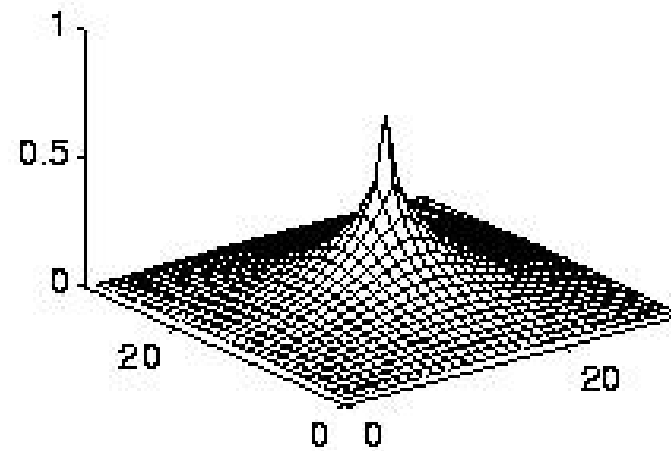
# Multigrid Motivation

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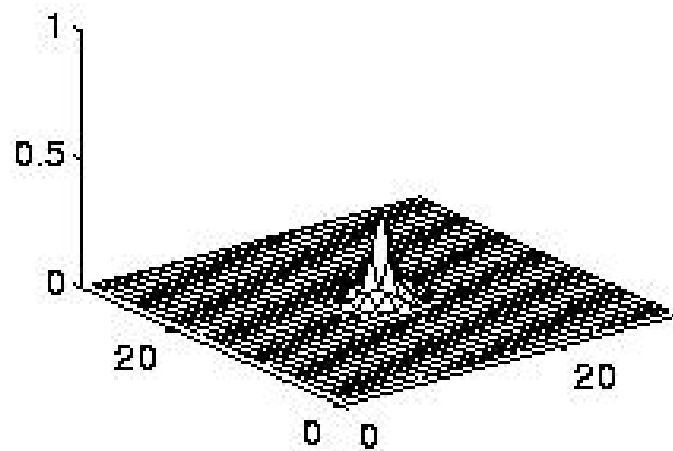
Right Hand Side



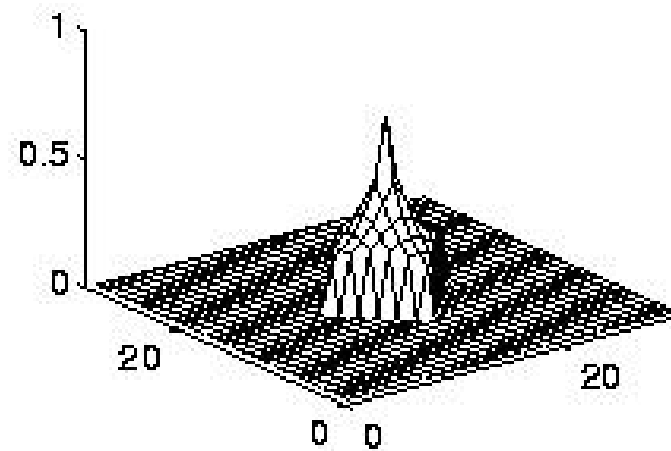
True Solution



5 steps of Jacobi



Best 5 step solution





# Multigrid Overview

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## ◦ Basic Algorithm:

- Replace problem on fine grid by an approximation on a coarser grid
- Solve the coarse grid problem approximately, and use the solution as a starting guess for the fine-grid problem, which is then iteratively updated
- Solve the coarse grid problem **recursively**, i.e. by using a still coarser grid approximation, etc.

## ◦ Success depends on coarse grid solution being a good approximation to the fine grid

## **Multigrid uses Divide-and-Conquer in 2 Ways**

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### **◦ First way:**

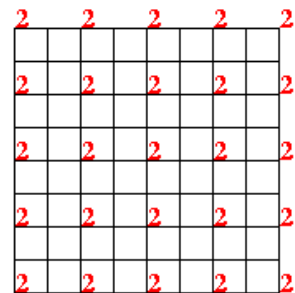
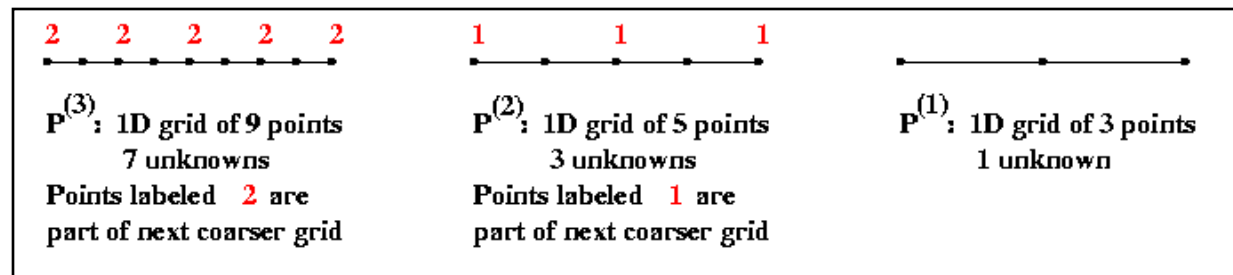
- Solve problem on a given grid by calling Multigrid on a coarse approximation to get a good guess to refine

### **◦ Second way:**

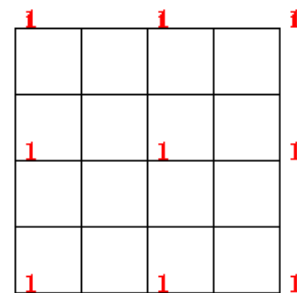
- Think of error as a sum of sine curves of different frequencies
- Same idea as FFT solution, but not explicit in algorithm
- Each call to Multigrid responsible for suppressing coefficients of sine curves of the lower half of the frequencies in the error (pictures later)

## Multigrid Sketch (1D and 2D)

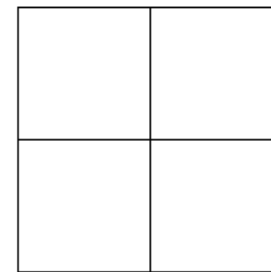
- Consider a  $2^m+1$  grid in 1D ( $2^m+1$  by  $2^m+1$  grid in 2D) for simplicity
- Let  $P^{(i)}$  be the problem of solving the discrete Poisson equation on a  $2^i+1$  grid in 1D ( $2^i+1$  by  $2^i+1$  grid in 2D)
  - Write linear system as  $T(i) * x(i) = b(i)$
- $P^{(m)}, P^{(m-1)}, \dots, P^{(1)}$  is sequence of problems from finest to coarsest



$P^{(3)}$ : 9 by 9 grid of points  
7 by 7 grid of unknowns  
Points labeled 2 are part of next coarser grid



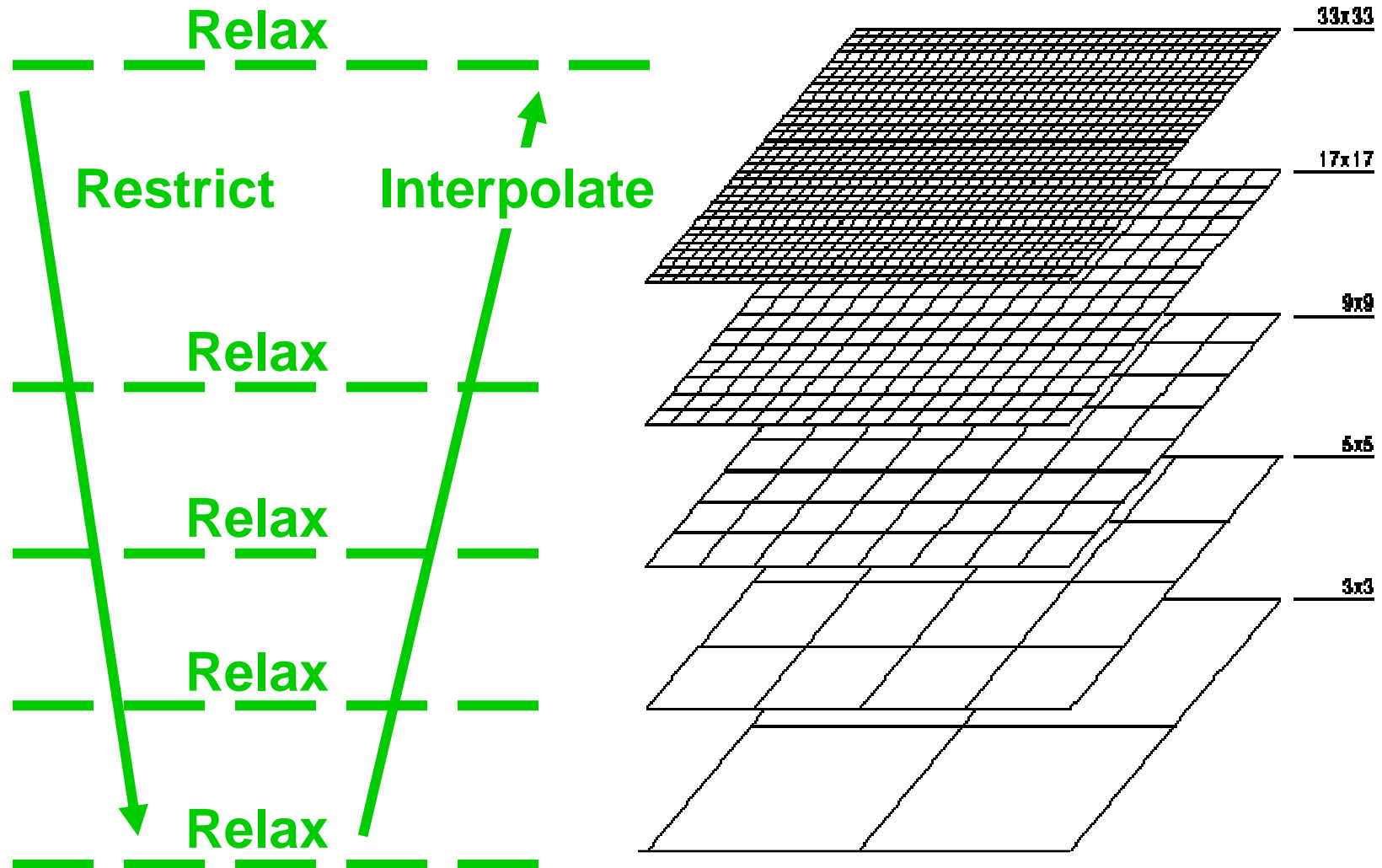
$P^{(2)}$ : 5 by 5 grid of points  
3 by 3 grid of unknowns  
Points labeled 1 are part of next coarser grid



$P^{(1)}$ : 3 by 3 grid of points  
1 by 1 grid of unknowns

# Multigrid Hierarchy

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Slide source: Geoffrey Fox

# Multigrid Operators

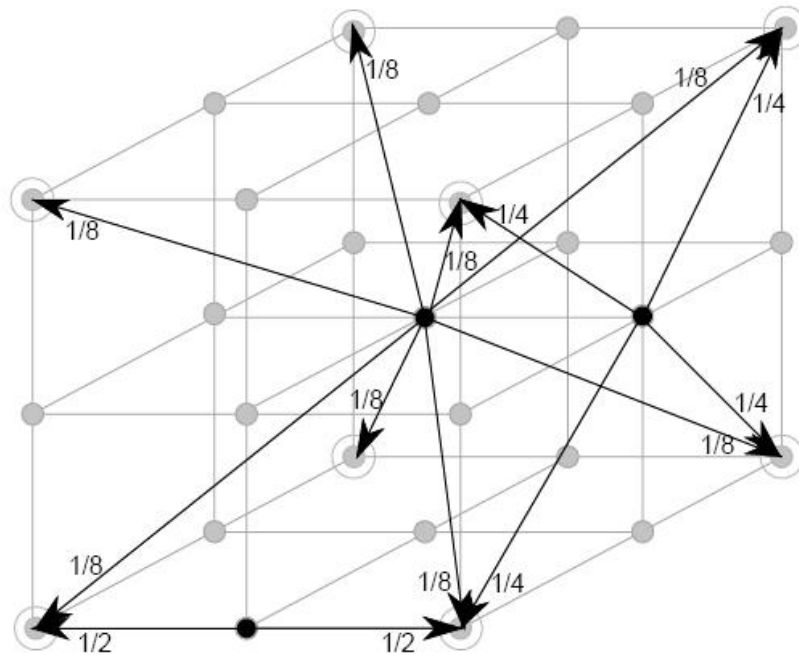
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- For problem  $P^{(i)}$  :
  - $b(i)$  is the RHS and
  - $x(i)$  is the current estimated solution
  - ( $T(i)$  is implicit in the operators below.)

} both live on grids of size  $2^i-1$
- All the following operators just average values on neighboring grid points
  - Neighboring grid points on coarse problems are far away in fine problems, so **information moves quickly** on coarse problems
- The **restriction operator  $R(i)$**  maps  $P^{(i)}$  to  $P^{(i-1)}$ 
  - Restricts problem on fine grid  $P^{(i)}$  to coarse grid  $P^{(i-1)}$  by sampling or averaging
  - $b(i-1) = R(i)(b(i))$
- The **interpolation operator  $ln(i-1)$**  maps an approximate solution  $x(i-1)$  to an  $x(i)$ 
  - Interpolates solution on coarse grid  $P^{(i-1)}$  to fine grid  $P^{(i)}$
  - $x(i) = ln(i-1)(x(i-1))$
- The **solution operator  $S(i)$**  takes  $P^{(i)}$  and computes an improved solution  $x(i)$  on same grid
  - Uses “weighted” Jacobi or SOR
  - $x_{\text{improved}}(i) = S(i)(b(i), x(i))$
- Details of these operators after describing overall algorithm

# Trilinear interpolation between the grids

## 3-D Prolongation

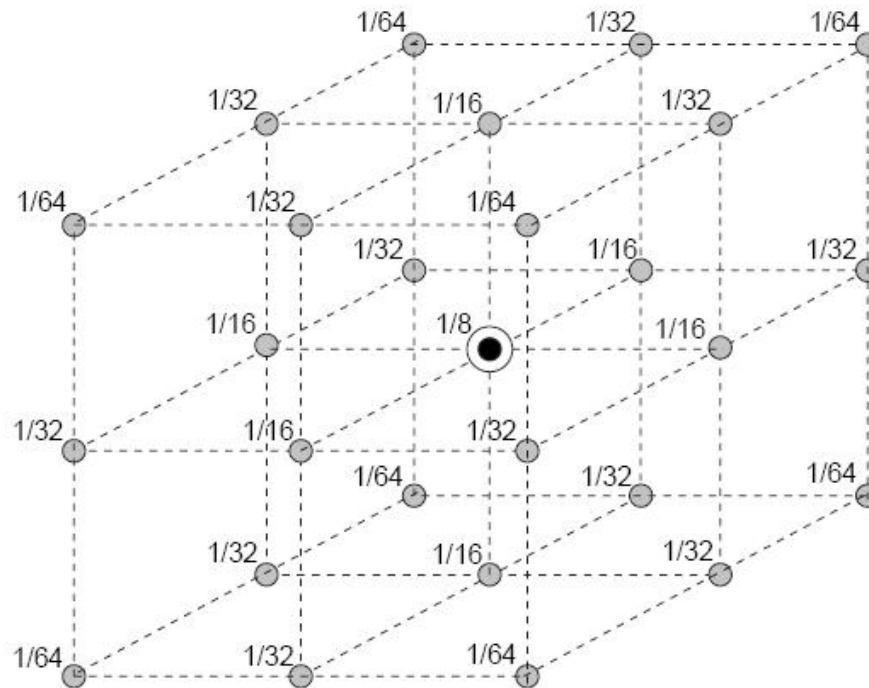


## 2-D Prolongation

$$\begin{bmatrix} 1 & 1 & 1 \\ \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \\ 1 & 1 & 1 \\ \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \end{bmatrix}$$

The arrows denote the coarse grid points to be used for interpolating the dense grid point. The numbers attached to the arrows denote the contribution of the specific coarse grid point.

## 3-D Restriction



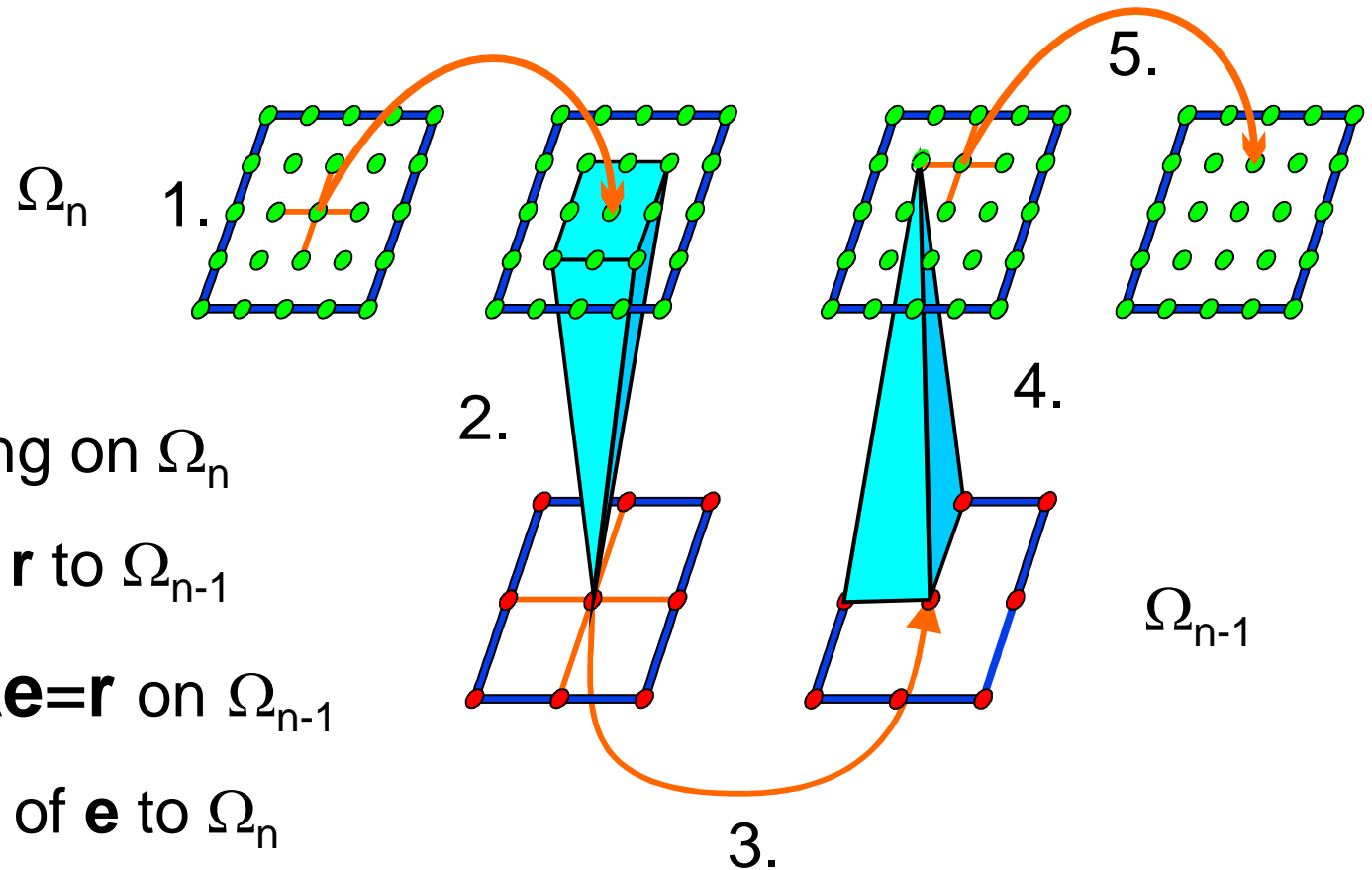
## 2-D Restriction

$$\begin{bmatrix} \frac{1}{16} & \frac{1}{8} & \frac{1}{16} \\ \frac{1}{8} & \frac{1}{4} & \frac{1}{8} \\ \frac{1}{16} & \frac{1}{8} & \frac{1}{16} \end{bmatrix}$$

A 27-point full weighting scheme is used. The number in front of each grid point denotes its weight in this operation.

## Multigrid method:

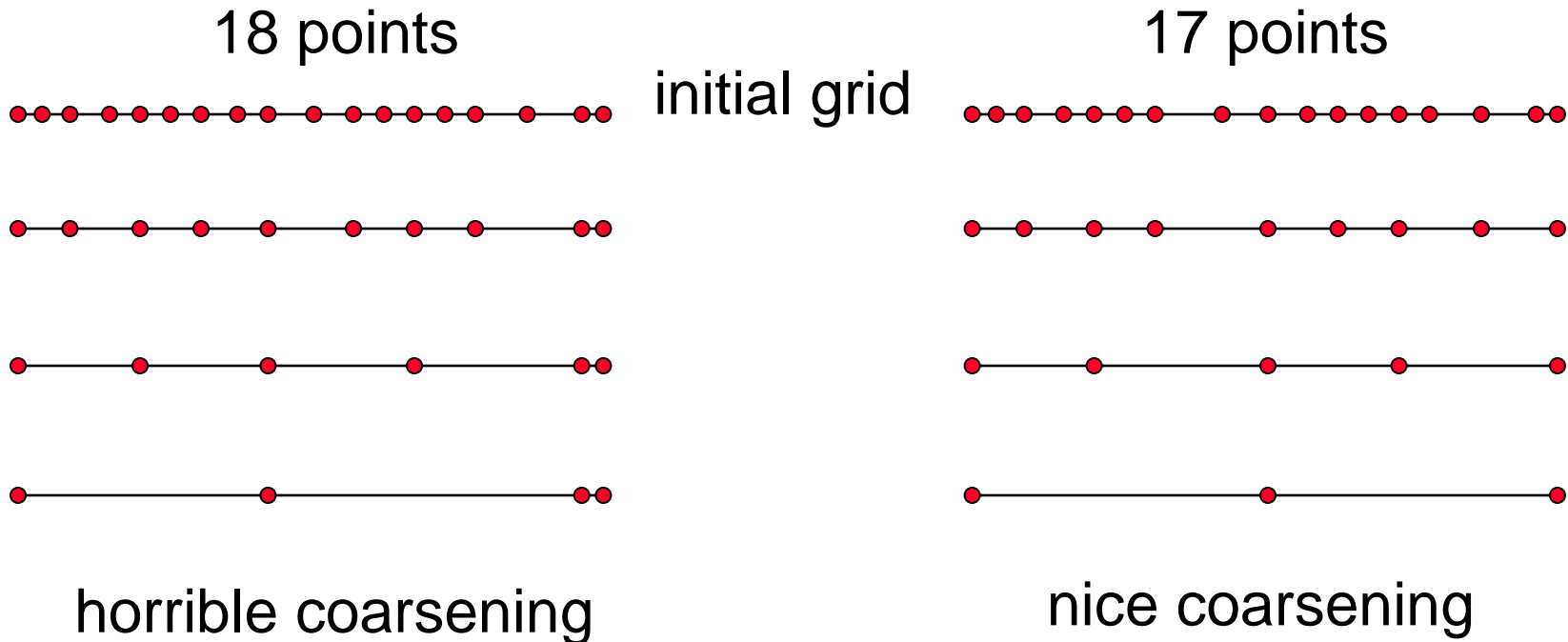
*two-grid iteration to solve  $A\mathbf{v}=\mathbf{f}$  on grid  $\Omega_n$*



1. pre-smoothing on  $\Omega_n$
2. restriction of  $\mathbf{r}$  to  $\Omega_{n-1}$
3. solution of  $A\mathbf{e}=\mathbf{r}$  on  $\Omega_{n-1}$
4. prolongation of  $\mathbf{e}$  to  $\Omega_n$
5. post-smoothing on  $\Omega_n$



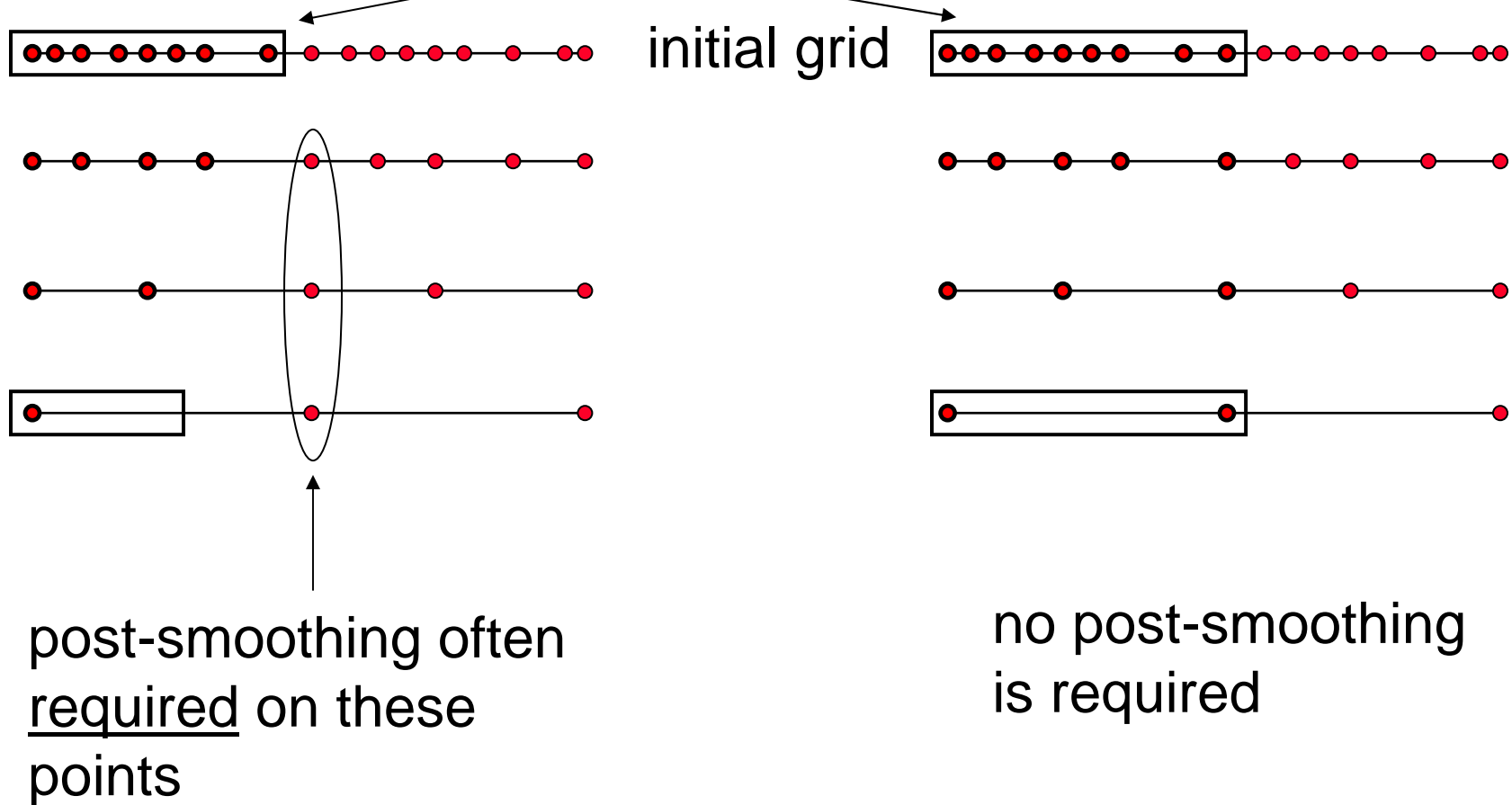
# Coarsening techniques: importance of boundaries I



When the number of points in one dimension is  $2^N+2$  ( $N$  being a natural number), a geometric mismatch is generated in the coarser grids, which show pronounced in-homogeneity. The convergence of the method is severely slowed down in these cases.

# Coarsening techniques: importance of boundaries II

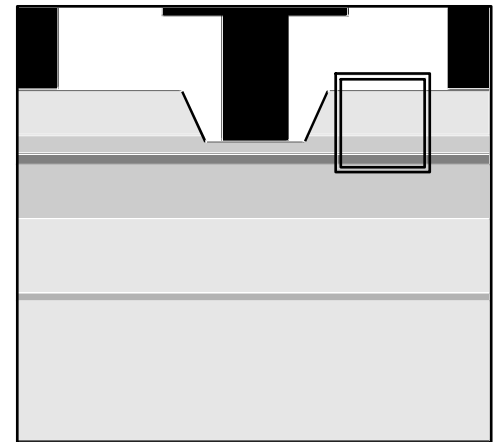
propagation of a contact (Neumann BC) through grids



# Coarsening techniques: simple application

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Grid	points	X-points	Y-points	coarsening
$\Omega_7$	22962	258	89	X,Y
$\Omega_6$	5850	130	45	X,Y
$\Omega_5$	1518	66	23	X,Y
$\Omega_4$	408	34	12	X,Y
$\Omega_3$	126	18	7	X,Y
$\Omega_2$	40	10	4	X
$\Omega_1$	24	6	4	X
$\Omega_0$	16	4	4	-



# Multigrid V-Cycle Algorithm

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Function MGV (  $b(i)$ ,  $x(i)$  )

... Solve  $T(i)*x(i) = b(i)$  given  $b(i)$  and an initial guess for  $x(i)$

... return an improved  $x(i)$

if ( $i = 1$ )

    compute exact solution  $x(1)$  of  $P^{(1)}$

only 1 unknown

    return  $x(1)$

else

$x(i) = S(i) (b(i), x(i))$

improve solution by

damping high frequency error

$r(i) = T(i)*x(i) - b(i)$

compute residual

$d(i) = I_{n(i-1)} ( \text{MGV}( R(i) ( r(i) ), 0 ) )$

solve  $T(i)*d(i) = r(i)$  recursively

$x(i) = x(i) - d(i)$

correct fine grid solution

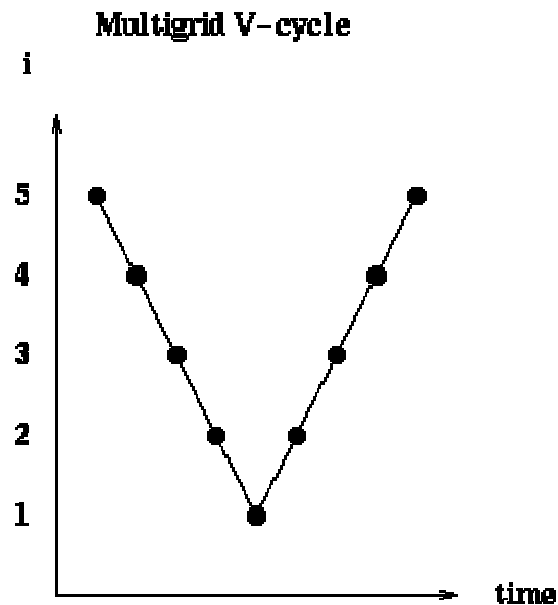
$x(i) = S(i) ( b(i), x(i) )$

improve solution again

    return  $x(i)$

## Why is this called a V-Cycle?

- ° Just a picture of the call graph
- ° In time a V-cycle looks like the following



## **Complexity of a V-Cycle on a Serial Machine**

- Work at each point in a V-cycle is  $O(\# \text{ unknowns})$
- Cost of Level  $i$  is  $(2^i - 1)^2 = O(4^i)$
- If finest grid level is  $m$ , total time is:

$$\sum_{i=1}^m O(4^i) = O(4^{m+1}) = O(\# \text{ unknowns})$$

# Full Multigrid (FMG)

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## ◦ Intuition:

- improve solution by doing multiple V-cycles
- avoid expensive fine-grid (high frequency) cycles

**Function FMG (b(m), x(m))**

**... return improved x(m) given initial guess**

**compute the exact solution x(1) of P(1)**

**for i=2 to m**

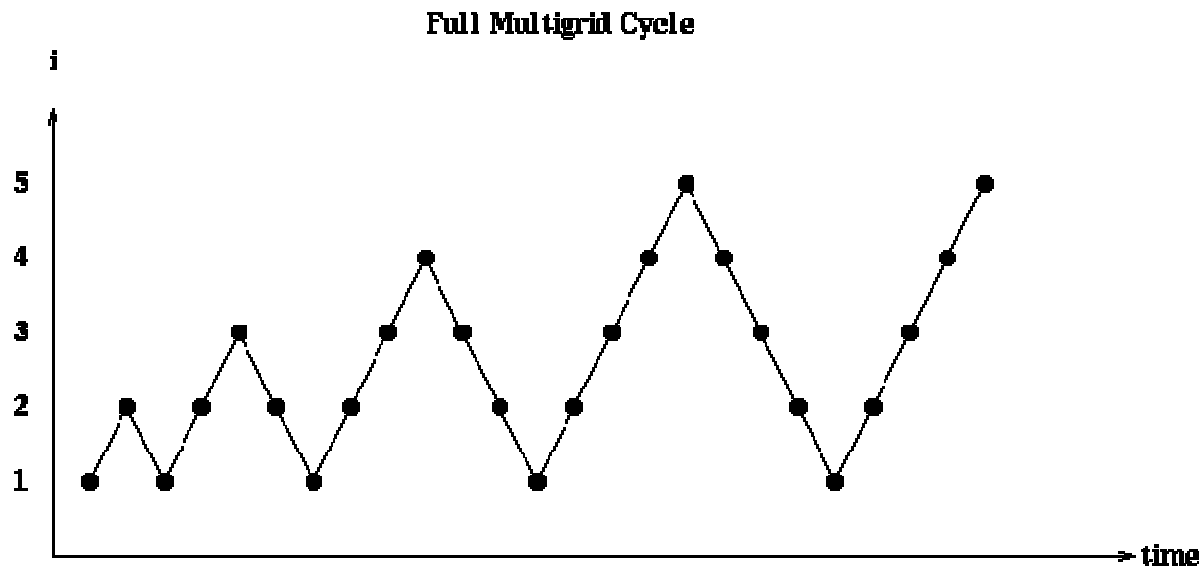
**$x(i) = \text{MGV} ( b(i), \text{In } (i-1) (x(i-1)) )$**

## ◦ In other words:

- Solve the problem with 1 unknown
- Given a solution to the coarser problem,  $P^{(i-1)}$ , map it to starting guess for  $P^{(i)}$
- Solve the finer problem using the Multigrid V-cycle

# Full Multigrid Cost Analysis

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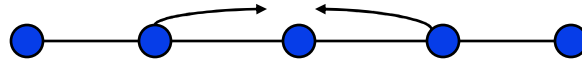
- One V for each call to FMG
  - people also use Ws and other compositions
- Serial time:  $\sum_{i=1}^m O(4^i) = O(4^m) = O(\# \text{ unknowns})$



## The Solution Operator $S(i)$ - Details

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- The solution operator,  $S(i)$ , is a weighted Jacobi
- Consider the 1D problem



- At level  $i$ , pure Jacobi replaces:

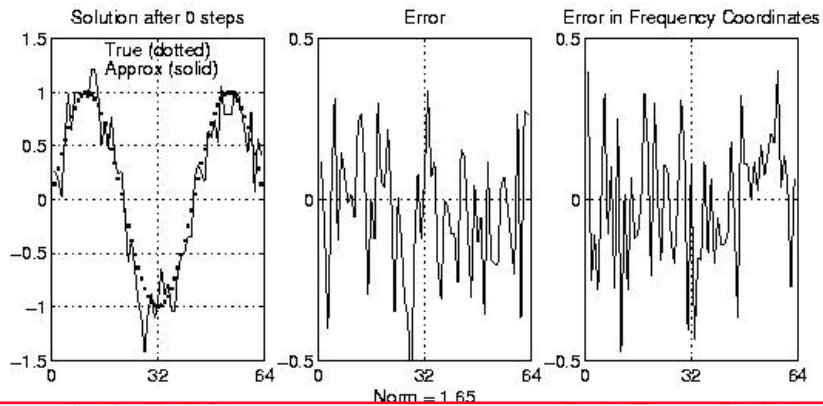
$$x(j) := 1/2 (x(j-1) + x(j+1) + b(j))$$

- Weighted Jacobi uses:

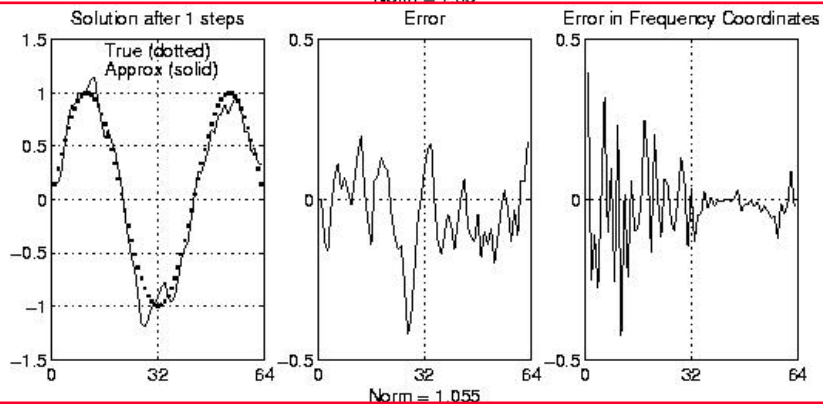
$$x(j) := 1/3 (x(j-1) + x(j) + x(j+1) + b(j))$$

- In 2D, similar average of nearest neighbors

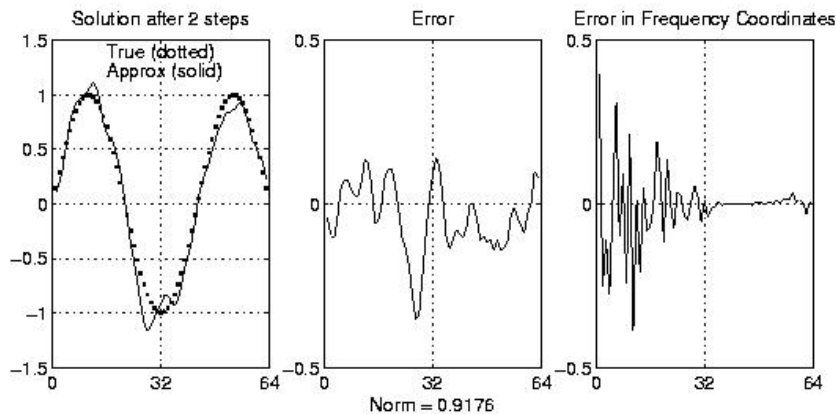
# Weighted Jacobi chosen to damp high frequency error



**Initial error**  
**“Rough”**  
**Lots of high frequency components**  
**Norm = 1.65**



**Error after 1 Jacobi step**  
**“Smoother”**  
**Less high frequency component**  
**Norm = 1.055**



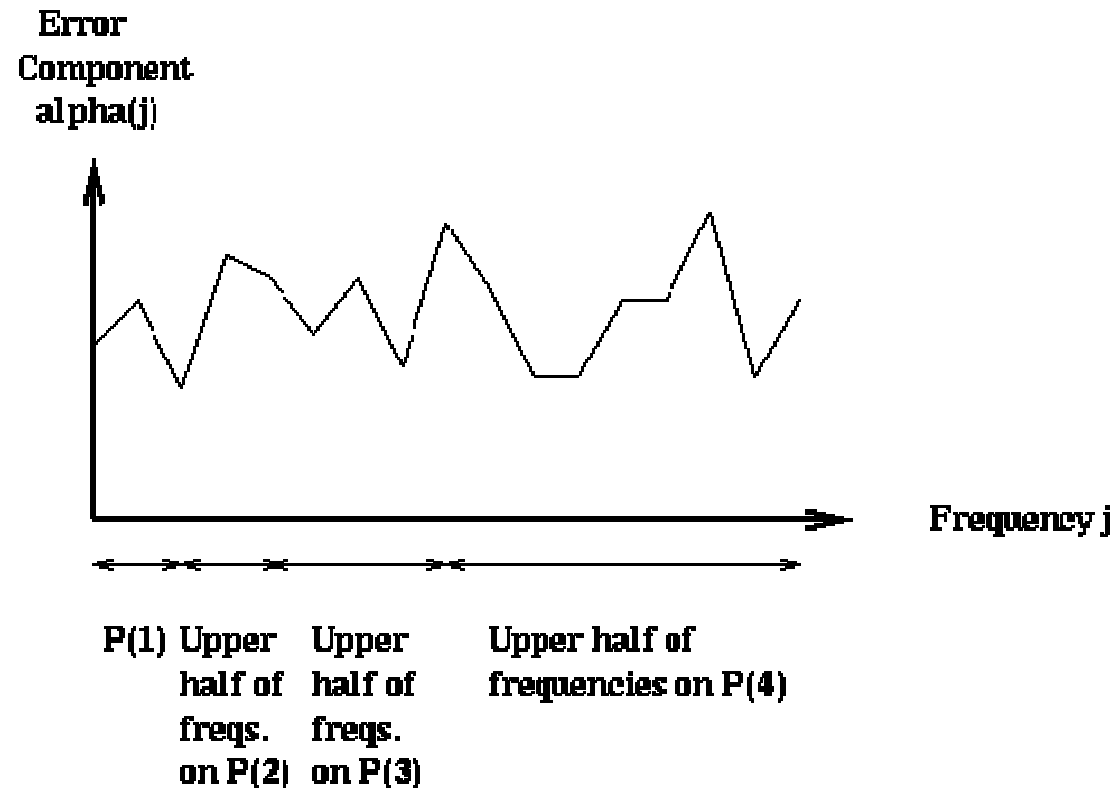
**Error after 2 Jacobi steps**  
**“Smooth”**  
**Little high frequency component**  
**Norm = .9176,**  
**won't decrease much more**

# Multigrid as Divide and Conquer Algorithm

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- ° Each level in a V-Cycle reduces the error in one part of the frequency domain

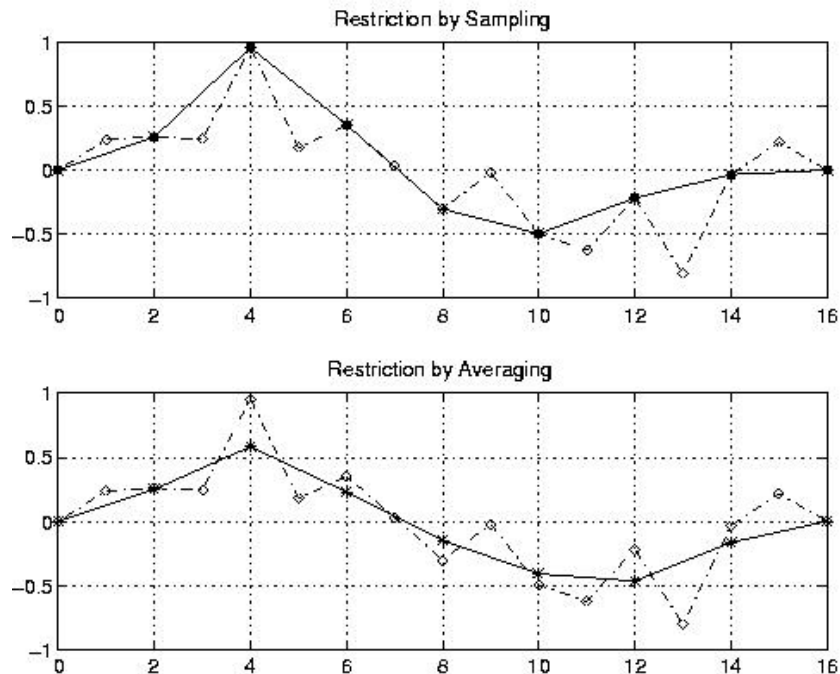
Schematic Description of Multigrid



# The Restriction Operator $R(i)$ - Details

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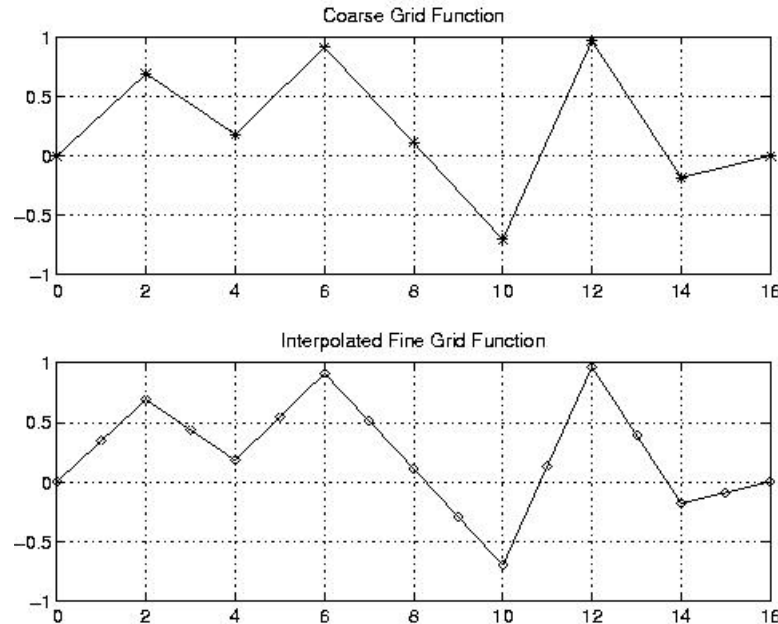
- The restriction operator,  $R(i)$ , takes
  - a problem  $P^{(i)}$  with RHS  $b(i)$  and
  - maps it to a coarser problem  $P^{(i-1)}$  with RHS  $b(i-1)$
- In 1D, average values of neighbors
  - $x_{\text{coarse}}(i) = 1/4 * x_{\text{fine}}(i-1) + 1/2 * x_{\text{fine}}(i) + 1/4 * x_{\text{fine}}(i+1)$



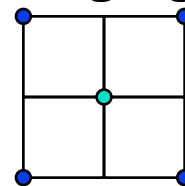
- In 2D, average with all 8 neighbors (N,S,E,W,NE,NW,SE,SW)

# Interpolation Operator

- The interpolation operator  $I_n(i-1)$ , takes a function on a coarse grid  $P^{(i-1)}$ , and produces a function on a fine grid  $P^{(i)}$
- In 1D, linearly interpolate nearest coarse neighbors
  - $x_{\text{fine}}(i) = x_{\text{coarse}}(i)$  if the fine grid point  $i$  is also a coarse one, else
  - $x_{\text{fine}}(i) = 1/2 * x_{\text{coarse}}(\text{left of } i) + 1/2 * x_{\text{coarse}}(\text{right of } i)$

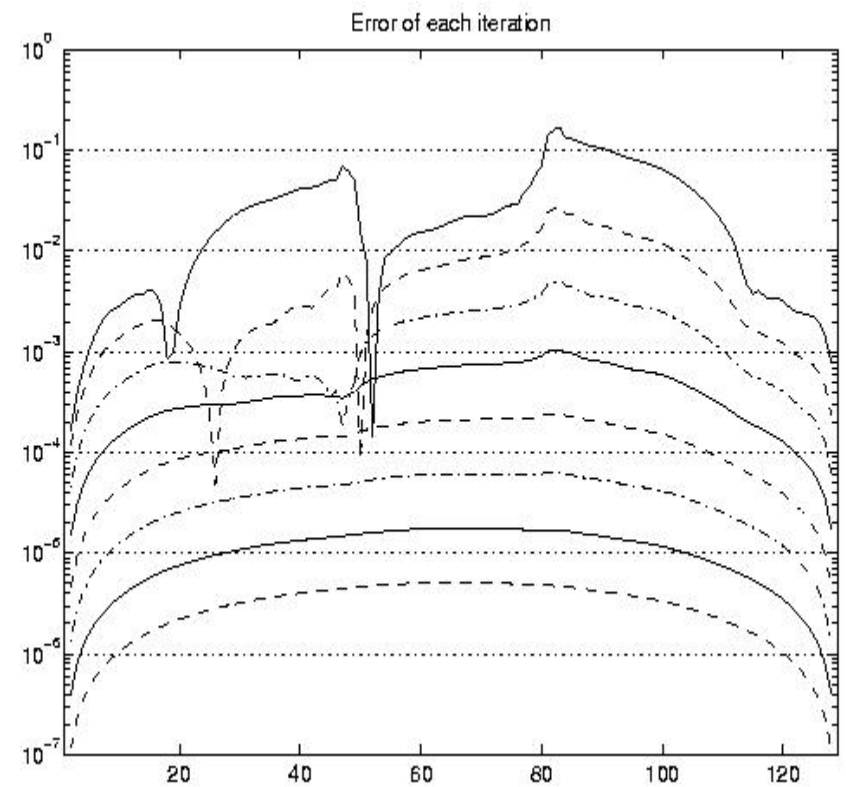
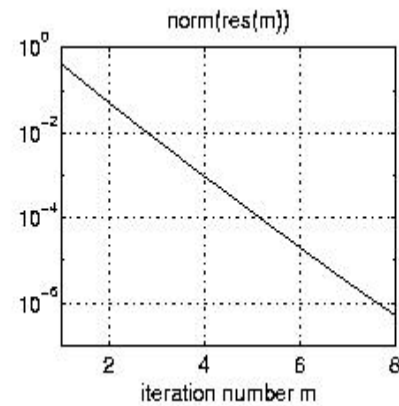
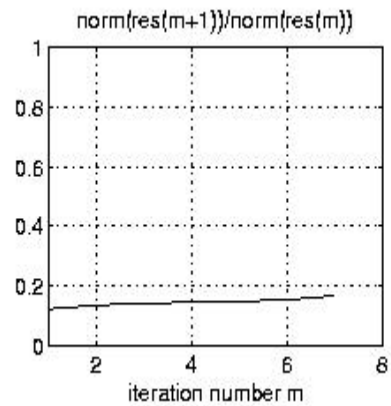
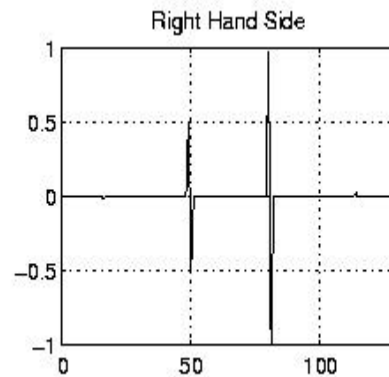
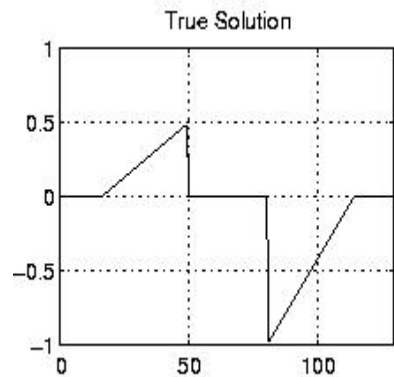


- In 2D, interpolation requires averaging with 2 or 4 nearest neighbors (NW,SW,NE,SE)



# Convergence Picture of Multigrid in 1D

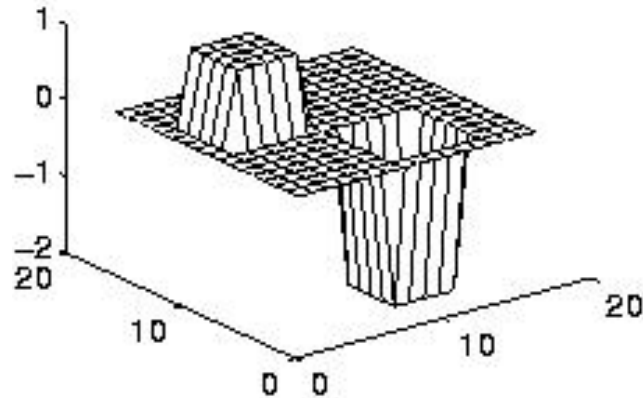
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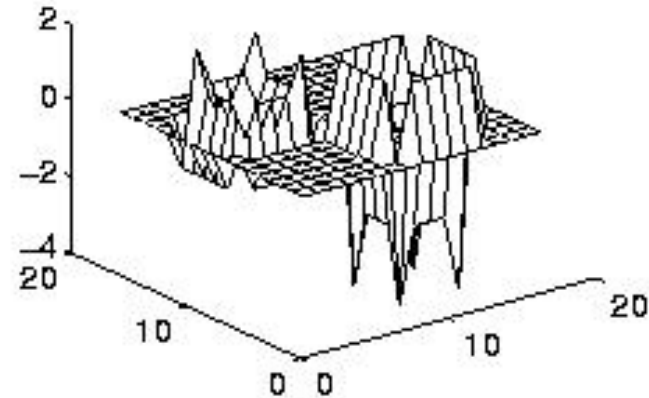
# Convergence Picture of Multigrid in 2D

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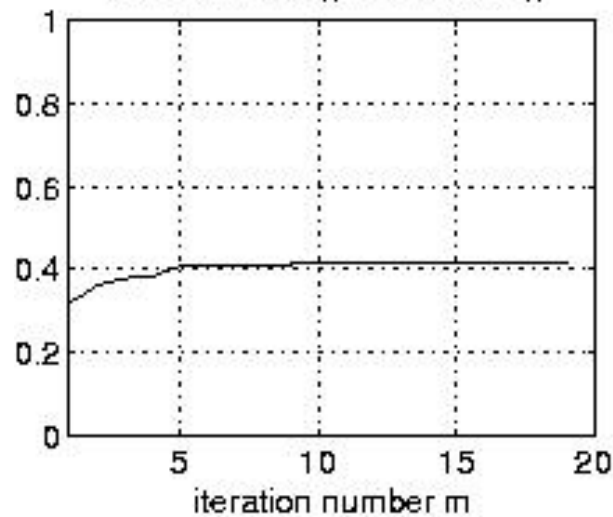
True Solution



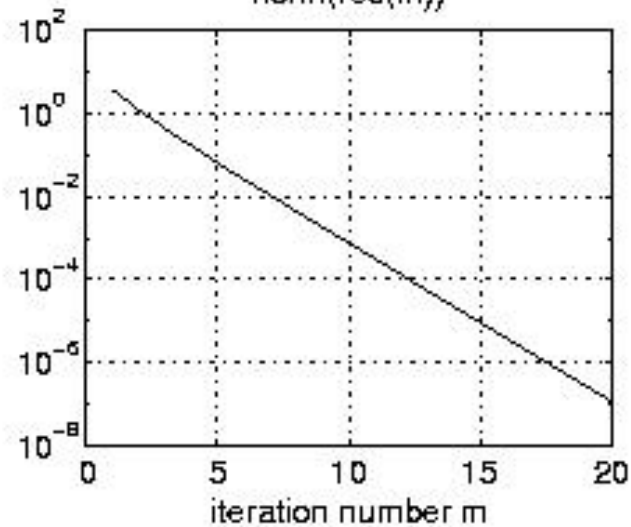
Right Hand Side



$\text{norm}(\text{res}(m+1))/\text{norm}(\text{res}(m))$



$\text{norm}(\text{res}(m))$



# Equilibrium Simulation Results

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PN diode	Grid points
<b>Grid level 6</b>	<b>129x129 x129</b>
<b>5</b>	<b>65 x 65 x 65</b>
<b>4</b>	<b>33 x 33 x 33</b>
<b>3</b>	<b>17 x17 x 17</b>
<b>2</b>	<b>9 x 9 x 9</b>
<b>1</b>	<b>5 x 5 x 5</b>

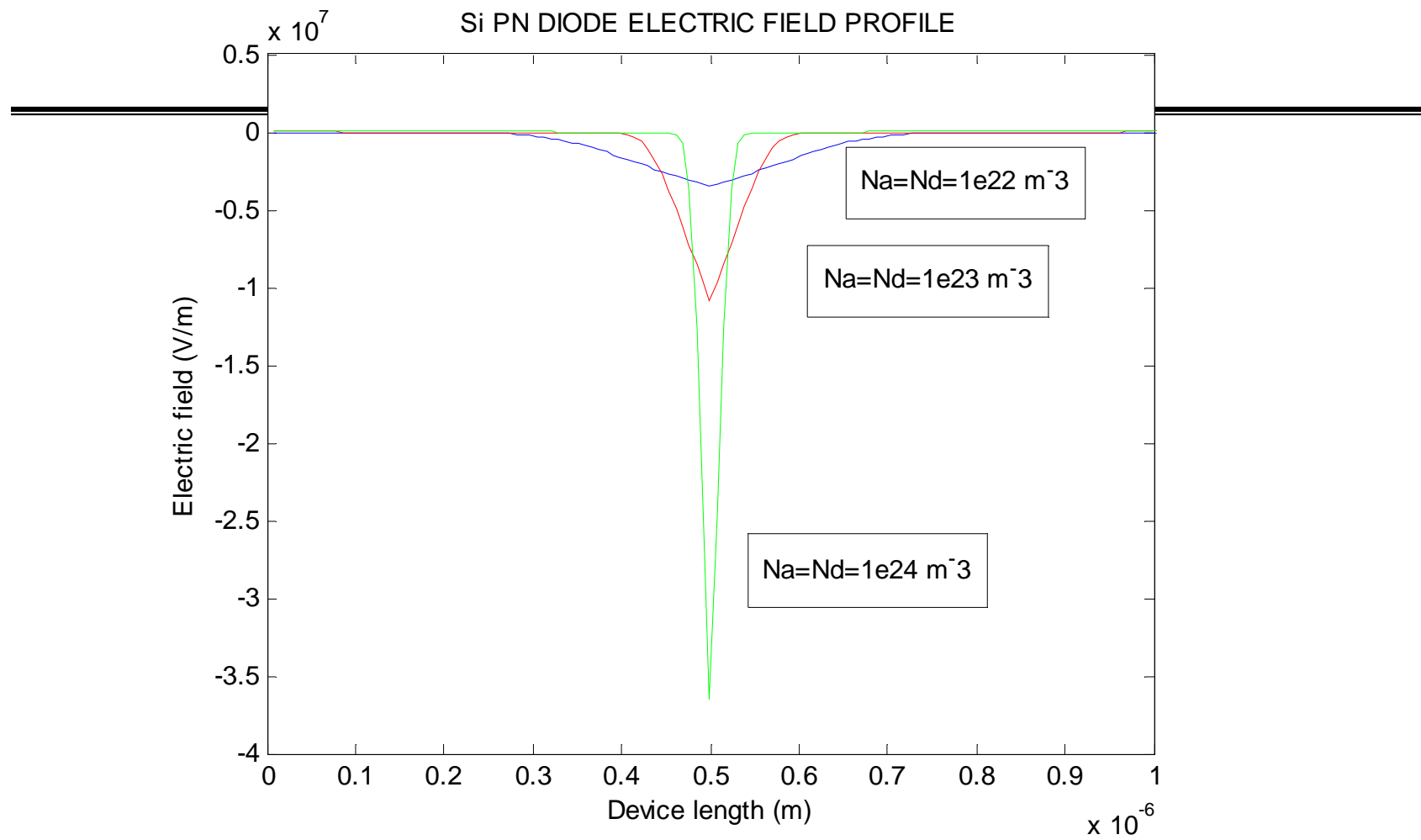
**Na (m<sup>-3</sup>)                      1.00E+22                      1.00E+23                      1.00E+24**

**Nd (m<sup>-3</sup>)                      1.00E+22                      1.00E+23                      1.00E+24**

**Depletion  
Width (m)                      0.601E-06                      0.206E-06                      0.069E-06**

Uniform grid organization on six levels (top panel). Comparison of depletion widths for different doping (bottom panel).





Na(m <sup>-3</sup> )	1.00E+22	1.00E+23	1.00E+24
Nd(m <sup>-3</sup> )	1.00E+22	1.00E+23	1.00E+24
Electric field(V/m)	-3.269E6	-1.1186E7	-3.7866E7